# Finitely Presented expansions of semigroups, algebras, and groups

Bakh Khoussainov

joint work with Denis Hirschfeldt (University of Chicago) and Alexei Miasnikov (Stevens Institute)

May 14, 2013

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- Finitely presented universal algebras
- The Bergstra-Tucker problem
- The NFP theorem
- Applications of the NFP theorem

## Definition

A **universal algebra** A is a tuple (A;  $f_1, \ldots, f_n, c_1, \ldots, c_m$ ), where

- $A \neq \emptyset$  is the *domain* of  $\mathcal{A}$ ,
- Each  $f_i$  is a total function  $A^{k_i} \to A$ , a *basic operation* of A.
- Each  $c_i$  is a distinguished element of A.

The **signature** of A is the sequence  $f_1, \ldots, f_n, c_1, \ldots, c_m$ .

### Definition

Each constant symbol and a variable *x* is a **term**. If  $t_1, \ldots, t_k$  are terms and *f* is function symbol of arity *k* then the expression  $f_i(t_1, \ldots, t_k)$  is a **term**.

Examples of terms are:

• f(f(x, y), z), f(x, f(y, z)), f(x, g(x)), f(f(x, a), f(a, a)), and

2 a, f(a, b), f(a, f(b, g(a))), f(g(a), a).

Terms in Example (2) are called ground terms.

The term algebra  $T_G$  is defined as follows:

- The domain is  $T_G$ , the set of all ground terms.
- For a basic operation symbol *f* of arity *k*, its value on the *k*-tuple (*t*<sub>1</sub>,..., *t<sub>k</sub>*) of ground terms is the ground term

 $f(t_1,\ldots,t_k).$ 

# The word problem

Let  $\ensuremath{\mathcal{A}}$  be a universal algebra generated by the constants. Then there exists a homomorphism

$$h: \mathcal{T}_G \to \mathcal{A}.$$

#### Definition

The word problem for  $\mathcal{A}$  is the set

$$WP(\mathcal{A}) = \{(p,q) \mid h(p) = h(q)\}.$$

#### Definition

The universal algebra A is **computably enumerable** if WP(A) is a computably enumerable set.

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# Definition

An equation is of the form t = t', where t and t' are terms. An equational specification S is a finite set of equations.

Examples:

• GroupAxioms, RingAxioms, BooleanAlgebraAxioms.

# Finitely presented universal algebras

Let *S* be an equational specification. Consider:

 $E(S) = \{(t, p) \mid t, p \in T_G \text{ and the equality } t = p \text{ can be}$ deduced from the equations in  $S\}.$ 

Define:

$$\mathcal{T}(S) = \mathcal{T}_G/E(S).$$

#### Definition

A universal algebra  $\mathcal{A}$  is **finitely presented** if there is an S such that  $\mathcal{A}$  is isomorphic to  $\mathcal{T}(S)$ .

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### Definition

A universal algebra A is **finitely presented** if there is an *S* such that A is isomorphic to T(S).

A computability-theoretic property:

•  $\mathcal{T}(S)$  is a c.e. algebra.

Algebraic and logical properties:

- T(S) is finitely generated and satisfies *S*.
- If A ⊨ S and A is finitely generated by the constants then A is a homomorphic image of T(S).
- $\mathcal{T}(S)$  is unique.

# Examples:

- The universal algebra ( $\omega$ ; S,  $2^x$ , 0) is not finitely presented.
- (Baumslag) The restricted wreath product of any non-trivial group G by Z is not finitely presented. In particular, the restricted wreath product of Z<sub>2</sub> by Z is not finitely presented (The lamplighter group).

# Making ( $\omega$ ; S, 2<sup>x</sup>, 0) finitely presented

Consider the following **expansion** of  $(\omega; S, 2^{\chi}, 0)$ :

 $(\omega; S, 2^{\chi}, 0, +, \times).$ 

This expansion is finitely presented. The presentation is this:

$$x + 0 = x, \ x + S(y) = S(x + y)$$

 $x \times 0 = 0, \ x \times S(0) = x, \ x \times S(y) = x \times y + x,$ 

$$2^{0} = S(0), \ 2^{x+S(0)} = 2^{x} \times S(S(0)).$$

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### Definition

An **expansion** of  $\mathcal{A} = (A; f_1, \dots, f_n)$  is  $\mathcal{A}' = (A; f_1, \dots, f_n, h_1, \dots, h_k)$ , where all *h* are new operations.

### Theorem (Bergstra-Tucker, 1980)

Every computably enumerable universal algebra with decidable word problem has finitely presented expansion. In the 1980s Bergstra and Tucker, and independently Goncharov, posed the following question:

Does every finitely generated computably enumerable universal algebra possess a finitely presented expansion?

Theorem (Kassymov (1987), Khoussainov (1994))

There exist finitely generated c.e. universal algebras that fail to possess finitely presented expansions.

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Do there exist standard finitely generated algebraic structures that fail to posses finitely presented expansions?

More precisely, do there exist finitely generated c.e. semigroups, groups, and algebras without finitely presented expansions?

# Effectively infinite universal algebras

### Definition

A universal algebra A is **effectively infinite** if there exists an algorithm that enumerates a sequence  $t_0, t_1, \ldots$  of ground terms such that they all represent pairwise distinct elements of A.

# Non-effectively infinite universal algebras

If  $\mathcal{A}$  is not effectively infinite then:

- Every expansion of A is not effectively infinite.
- ② All infinite sub-algebras of  $\mathcal{A}$  are not effectively infinite.
- So For every term t(x) and element *a* of A, the sequence

 $a, t(a), t(t(a)), \ldots$ 

is eventually periodic. In particular if a group is not effectively infinite then it is a periodic group.

All infinite homomorphic images of A are not effectively infinite.

# Non-effectively infinite semigroups

# Theorem (Khoussainov, Hirschfeldt)

There exists a c.e., f.g. and non effectively infinite semigroup.

**Proof.** Consider  $\mathcal{A} = (\{x, y\}^*; \cdot)$ . Take  $Z \subseteq \{x, y\}^*$ , and define:

 $R(Z) = \{u \mid u \text{ contains a subword from } Z\}.$ 

The relation  $\eta(Z)$ 

 $(u, v) \in \eta(Z) \leftrightarrow u = v \lor (u \in R(Z) \& v \in R(Z)).$ 

is a congruence relation on  $\mathcal{A}$ . Set:

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#### Lemma

There exists a simple subset  $Z \subseteq \{x, y\}^*$  such that

• The semigroup  $A(Z) = A/\eta(Z)$  is infinite.

2 For each *n*,  $r_n$  the number of strings in *Z* of length *n* is bounded by *n* and  $r_0 = r_1 = \ldots = r_5 = 0$ .

The lemma is proved by a careful Post type of construction of a simple set. We use Joe Miller's lemma.

This lemma proves the theorem.

### Definition

A universal algebra  $\mathcal{A}$  from a class K is **residually finite** if for any two distinct elements a and b of  $\mathcal{A}$ , there exists a homomorphism of  $\mathcal{A}$  onto a finite universal algebra from the class K such that the images of a and b under the homomorphism are distinct.

#### Lemma

Let  $\mathcal{A} = \mathcal{T}_G / E$  be a residually finite universal algebra. For all distinct elements x and y of  $\mathcal{A}$  there exists a set  $S(x, y) \subseteq T_G$  such that each of the following is satisfied:

• S(x, y) is a computable set.

2 
$$x \in S(x, y)$$
 and  $y \notin S(x, y)$ .

The set S(x, y) is a union of E-equivalence classes.

# Non-effective infinity and residual finiteness

### Theorem (Khoussainov and Miasnikov)

Let  $\mathcal{A} = \mathcal{T}_G / E$  be a f.g. and c.e. universal algebra.

If A is residually finite and non-effectively infinite, then all c.e. expansions of A are residually finite.

In other words, residual finiteness is invariant under expansions if non-effective infinity is assumed.

**Proof**. Let  $\mathcal{A}'$  be a c.e. expansion of  $\mathcal{A}$ .

- Set *C*<sub>(*a,b*)</sub> be the minimal congruence relation generated by the pair (*a*, *b*).
- There exists a procedure that given (a, b) extracts an algorithm generating  $C_{(a,b)}$ .
- Select x ≠ y from A'. Want to show that we can distinguish them in a finite homomorphic image of A'.

# Ingredients of the proof II

Define  $E_{(x,y)}$ :

 $(a,b) \in E_{(x,y)} \text{ iff } \forall u \in S(x,y) \ \forall v \notin S(x,y)((u,v) \notin C_{(a,b)}).$ 

#### \_emma

The binary relation  $E_{(x,y)}$  is a congruence relation of  $\mathcal{A}'$ .

This is based on the following two lemmas (by Malcev):

- $(u, v) \in C_{(a,b)}$  iff  $\exists e_0, ..., e_n \in A$  and  $\exists t_0, ..., t_{n-1}$  such that  $u = e_0, v = e_n$ , and  $\{e_i, e_{i+1}\} = \{t_i(a), t_i(b)\}$  for  $i \le n 1$ .
- Is a congruence of  $\mathcal{A}$  iff all algebraic terms respect E.

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- E is a congruence of A iff all algebraic terms respect E.

#### Lemma

The relation  $E_{(x,y)}$  is a co-c.e.congruence relation.

The transversal of every co-c.e. equivalence relation is a c.e. set. Hence, the transversal must be a finite set.

Thus, E(x, y) has a finite index. Therefore  $\mathcal{A}'/E(x, y)$  is the desired homomorphic image.

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If A is an infinite, finitely generated, computably enumerable, non-effectively infinite and residually finite universal algebra then no expansion of A is finitely presented.

#### Proof.

Assume A' is a f.p. expansion of A. Then A' is residually finite as we have proved. The next lemma implies the theorem:

#### \_emma (Malcev)

If a finitely presented algebra is residually finite then the word problem for the algebra is decidable.

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### Lemma (Malcev)

If a finitely presented algebra is residually finite then the word problem for the algebra is decidable.

### Corollary

There exists a finitely generated, non-effectively infinite, residually finite and computably enumerable semigroup that fails to possess finitely presented expansions.

### Proof.

It suffices to note that the semigroup  $\mathcal{A}(Z)$  constructed above is residually finite.

Consider the algebra of polynomials over the field *k*:

$$\mathcal{F}=k\{x_1,\ldots,x_d\},\ d>1.$$

Represent  $\mathcal{F}$  as

$$\mathcal{F}_0 \oplus \mathcal{F}_1 \oplus \mathcal{F}_2 \oplus \ldots$$

Write each polynom p in  $\mathcal{F}$  as the sum of **homogeneous** polynomials:

$$p=p_1+p_2+\ldots+p_k,$$

where  $p_1 \in \mathcal{F}_{n_1}, \ldots, p_k \in F_{n_k}$  and  $n_1 < n_1 < \ldots < n_k$ .

# Homogeneous ideals

# Definition

A ideal is **homogeneous** if it is generated by homogeneous polynomials.

#### Lemma

An ideal I of  $\mathcal{F}$  is homogenous if and only if it contains homogeneous components of all of its members.

#### Lemma

If I is homogeneous then  $\mathcal{A} = \mathcal{F}/I$  can be represented as

$$\mathcal{A}_0 \oplus \mathcal{A}_1 \oplus \mathcal{A}_2 \oplus \ldots,$$

where  $A_n = (F_n + I)/I$  for all n.

Let *I* be an ideal generated by homogeneous polynomials:

 $h_1, h_2, h_3, \ldots, \text{ where } 2 \le deg(h_1) \le deg(h_2) \le \ldots$ 

Set:  $r_n = \#$ (polynomials of degree *n* in this sequence).

### Theorem (Golod-Shafarevich Theorem)

Assume that for some  $\epsilon$  such that  $0 < \epsilon \le d/2$  we have the following inequality:

$$r_n \leq \epsilon^2 (d-2\epsilon)^{n-2}$$
 for all  $n \geq 2$ .

Then the algebra A = F/I has an infinite dimension.

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Let  $W_0, W_1, \ldots$  be an enumeration of all c.e. subsets of  $\mathcal{F}$ . Uniformly on *i* proceed as follows. Wait for two polynomials *f* and *g* that appear in  $W_i$  such that:

- Both *f* and *g* are of the form  $f = f_1 + h_1$ ,  $g = g_1 + h_2$ , and  $f_1 = g_1$ .
- 2 The degrees of all homogeneous components occurring in  $h_1$  and  $h_2$  are greater than i + 10.

Put the homogeneous components present in  $h_1$  and  $h_2$  into H.

Set  $I = \langle H \rangle$  and consider the algebra

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$$\mathcal{A} = \mathcal{F}/I.$$

The algebra  $\mathcal{A} = \mathcal{F}/I$  is finitely generated, infinite, computably enumerable, non-effectively infinite, residually finite, and nil-algebra.

For the proof one needs to check that the conditions of the Golod-Shafarevich theorem are satisfied. Hence,

Theorem (Khoussainov and Miasnikov)

There exists a finitely generated, non-effectively infinite, residually finite and computably enumerable algebra that fails to possess finitely presented expansions.

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### Theorem (Khoussainov and Miasnikov)

There exists a finitely generated, non-effectively infinite, residually finite and computably enumerable algebra that fails to possess finitely presented expansions.

For every prime number p there exists a finitely generated, non-effectively infinite, residually finite and computably enumerable p-group that fails to possess finitely presented expansions.

**Proof**. Consider the muliplicative substructure *G* of A generated by 1 + x and 1 + y. One can show the following:

- G is infinite.
- $\bigcirc$  G is a group.
- G is residually finite.

The group G is non-effectively infinite finite. Apply the NFP theorem.

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There exists a finitely generated, residually finite, computably enumerable and effectively infinite algebra that fails to possess finitely presented expansions.

Let *I* be the homogenious ideal generated by set *Z* constructed for the semigroup. The algebra  $\mathcal{A} = \mathcal{F}/I$  is the desired algebra.

- Does there exist effectively infinite groups that fail to possess finitely presented expansions?
- Are there examples of c.e. finitely generated groups, semigroups, and algebras that fail to possess quasi-equational presentations in all expansions?
- Obes there exist a finitely presented universal algebra which is ML-random?
- Ooes there exist an effectively infinite and ML-random algebra?