

On quantum Kolmogorov complexity and their relationships

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Preface

- Why QKC?
 - It can be a new tool to analyze quantum state in information theoretical point of view (ex. Miyadera[2012])
 - It can be a basic notion to consider “quantum randomness”
- Today I will talk about
 - Two definition of QKC by Berthiaume et al. and Gács and some known facts about them,
 - My conjecture (quantum Levin’s coding theorem).

Axioms for quantum mechanics

- To every quantum system \mathcal{S} uniquely associated is a Hilbert space \mathbf{H} called the state space of \mathcal{S} . States of \mathcal{S} are represented by unit vectors $|\psi\rangle$ in \mathbf{H} .

Ex.) Qubit system ... \mathbb{C}^2 , $\{|0\rangle, |1\rangle\}$

- For every unit vector in \mathbf{H} , there exists a state represented by that vector.

Ex.) $\frac{|0\rangle + |1\rangle}{\sqrt{2}}$

Axioms for quantum mechanics

- The state space of the composite system $\mathcal{S} = \mathcal{S}_1 + \mathcal{S}_2$ consisting of the system \mathcal{S}_1 with the state space H_1 and of the system \mathcal{S}_2 with the state space H_2 is the tensor product $H_1 \otimes H_2$.

Ex.) the space of n-letter qubit strings ... $\{\mathbb{C}^2\}^{\otimes n}$

$$\{|0\dots 0\rangle, \dots, |1\dots 1\rangle\}$$

→ the space of all qubit strings ... $\bigoplus_{n=0}^{\infty} \{\mathbb{C}^2\}^{\otimes n}$

Axioms for quantum mechanics

- (discrete) Time evolution of quantum states is expressed by a unitary operator U :

$$|\psi\rangle \rightarrow U|\psi\rangle \rightarrow U^2|\psi\rangle \rightarrow \dots$$

Several definitions of QKC

- Use quantum Turing machine
 - Vitányi [2001]
 - Berthiaume et al. [2001]
- “Quantize” universal semimeasure
 - Gács [2001]

Quantum Turing machine (QTM)

- $M = (Q, \Sigma, \delta)$
- $\delta : Q \times \Sigma \times Q \times \Sigma \times \{L, R\} \rightarrow \tilde{\mathbb{C}}$
- δ uniquely defines an operator \mathbf{U} on \mathbf{H}_{qtm} :

$$U|q, T, \zeta\rangle = \sum_{\substack{(p, \tau, d) \\ \in Q \times \Sigma \times \{-1, 1\}}} \delta(q, T(\zeta), p, \tau, d) |p, T_{\zeta}^{\tau}, \zeta + d\rangle$$

- We call this machine QTM only if \mathbf{U} is unitary.

Existence of universal QTM

Theorem (Bernstein – Vazirani, 1997) *There exists a QTM M_0 such that*

$$\forall M \exists s_m \forall T \forall \delta > 0 \forall |\psi\rangle \quad \|M_0(s_m, T, \delta, |\psi\rangle) - M(|\psi\rangle)\|_{Tr} < \delta.$$

This machine simulates each QTM with polynomial slowdown.

Qubit complexity (Berthiaume et al. 2001)

- $QC^\delta(|\psi\rangle) = \min \left\{ l(|\varphi\rangle) \mid \|M_0(|\varphi\rangle) - |\psi\rangle\|_{Tr} < \delta \right\}$
- $QC^\downarrow(|\psi\rangle) = \min \left\{ l(|\varphi\rangle) \mid \forall k \|M_0(|\varphi\rangle, k) - |\psi\rangle\|_{Tr} < \frac{1}{k} \right\}$

Problem (invariance of QC)

- (Classical case) for universal TM M_0 and any TM M , there exists a constant c_M such that

$$C_M(w) \leq C_{M_0}(w) + c_M.$$

(Just input $\langle s_M, w \rangle$ to M_0)

- Bernstein – Vazirani UQTM requires the Halting time as an input ... can't show the invariance of qubit complexity in the same manner as above.

Strongly UQTM (Müller, 2007)

Theorem (Müller, 2007) *There exists a QTM M_1 such that*

$$\forall M \forall |\psi\rangle \exists |\psi_M\rangle \forall \delta > 0 \quad \|M_1(|\psi_M\rangle, \delta) - M(|\psi\rangle)\|_{Tr} < \delta$$

for each $|\psi\rangle$ for which $M(|\psi\rangle)$ is defined, where

$$l(|\psi_M\rangle) \leq l(|\psi\rangle) + c_M.$$

$$\rightarrow \left\{ \begin{array}{l} QC_{M_0}^\Delta(|\psi\rangle) \leq QC_M^\delta(|\psi\rangle) + c_{\Delta, \delta, M} \quad (\delta < \Delta) \\ QC_{M_1}^\downarrow(|\psi\rangle) \leq QC_M^\downarrow(|\psi\rangle) + c_M \end{array} \right.$$

Relation between classical complexity

Theorem (Müller, 2009)

$$QC^\downarrow(|x\rangle) = C(x) + O(1).$$

Question: Is the notion of QTM indispensable for the definition of qubit complexity?

Is there an equivalent definition of qubit complexity on the notion of classical Turing machine?

Density matrix

- $\rho \in \mathbb{M}_N$ is called *Density matrix* if

$$\rho \geq 0, \quad \text{Tr}(\rho) = 1.$$

$$\rightarrow \rho = \sum_{i=1}^N \lambda_i |\psi_i\rangle\langle\psi_i|, \quad \text{where} \quad \lambda_i \geq 0, \quad \sum_{i=1}^N \lambda_i = 1.$$

- Density matrix represents a probability distribution of quantum states,

$$\{\lambda_1, \dots, \lambda_N; |\psi_1\rangle, \dots, |\psi_N\rangle\}.$$

Universal semi-density matrix (Gács, 2001)

- Semi-density matrix ρ (that is, $\rho \in \mathbb{M}_N$ with $\rho \geq 0$, $\text{Tr}(\rho) \leq 1$) is *lower-semicomputable* if there is a sequence $\{\rho_n\}$ of elementary matrix such that

$$\rho_n \leq \rho_{n+1}, \rho_n \rightarrow \rho \ (n \rightarrow \infty).$$

- Lower-semicomputable semi-density matrix μ is *universal* if for any LSC-SDM ρ there is a constant $c_\rho > 0$ such that

$$\rho \leq c_\rho \mu.$$

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