

# An integral test for Schnorr randomness and its applications

Kenshi Miyabe

Research Institute for Mathematical Sciences, Kyoto University

Some recent researches show that some classical theorems with “almost everywhere” can be converted to effective version with “for each random point”. One example is that a real is Martin-Löf random iff each computable function of bounded variation is differentiable at the real. The author [1] has given a Kurtz randomness version where we find that an integral test is a useful tool to study the relation between algorithmic randomness and computable analysis. Here we give a version of Schnorr randomness. Consider a computable metric space  $(X, d, \alpha)$  and a computable measure  $\mu$  on it.

**Definition 1.** An integral test for Schnorr randomness is a nonnegative lower semicomputable function  $t : X \rightarrow \overline{\mathbb{R}}^+$  such that  $\int t d\mu$  is a computable real.

**Theorem 1.** A point  $z$  is Schnorr random iff  $t(z) < \infty$  for each integral test  $t$  for Schnorr randomness.

As an application we have another effectivized version of a classical theorem.

**Theorem 2.** Let  $f, g : \subseteq X \rightarrow \mathbb{R}$  be the differences between two integral tests for Schnorr randomness. Then  $f(x) = g(x)$  for each Schnorr random point iff  $\|f - g\|_1 = 0$ .

The class of the differences between two integral tests for Schnorr randomness has some characterizations and is an important class.

**Definition 2.** A function  $f : \subseteq X \rightarrow \mathbb{R}$  is  $L^1$ -computable with an effective code if there exists a computable sequence  $\{s_n\}$  of finite rational step functions such that  $f(x) = \lim_n s_n(x)$  and  $\|s_{n+1} - s_n\|_1 \leq 2^{-n}$  for all  $n$ .

**Definition 3.** A function  $f : \subseteq X \rightarrow \mathbb{R}$  is Schnorr layerwise computable if there exists a Schnorr test  $U_n$  such that the restriction  $f|_{X \setminus U_n}$  is uniformly computable.

**Theorem 3.** Let  $\text{SR}$  be the set of Schnorr random points. Then

$$\begin{aligned} & \{f|_{\text{SR}} \mid f \text{ is the difference between two integral tests for Schnorr randomness}\} \\ = & \{f|_{\text{SR}} \mid f \text{ is an } L^1\text{-computable function with an effective code}\} \\ = & \{f|_{\text{SR}} \mid f \text{ is Schnorr layerwise computable and its } L^1\text{-norm is computable}\} \end{aligned}$$

## References

1. Miyabe, K.: Characterization of Kurtz randomness by a differentiation theorem, submitted